

PIC Architecture

Instruction Set & Operations

PICs-Instruction Set

- Have Covered Instruction Set Basics
 - Accumulator Architecture
 - Direct addressing
 - Indirect addressing
- Now lets look at the instructions

MOVE instructions

- PIC spend a lot of time moving data around as data stored in memory
- `movlw 20`
 - Move the hex value H'20' into W
 - To load a decimal must use the correct assembler directive - D'20'
- `movlw -2`
 - Loads B'1111 1110 into WREG

More on MOVE

- Initialization of a variable
 - `movlw B'11100000'`
 - `movwf TRIST`
 - Initialize the PORTB data direction register
- Assembler MACRO
 - `MOVLF B'11100000',TRISB`
 - Assembled into the two instructions above

The movff instruction

- `movff PORTB, PORTB_COPY`
 - `movff` - a two-word instruction
 - Thus source and destination have 12-bit addresses
 - Source – instruction bits + BSR
 - Destination – instruction 2nd byte
 - Moves data from a 1-byte source to a 1-byte destination.
 - For instruction memory as efficient as the regular move instruction.

The movwf

- `movwf f(,Banked)` – Move WREG to f
- For storing back the result of an operation
- Does not affect status bits

The movf instruction

- Move the value and set the status register bits appropriately
- Affects N and Z bits

Move summary

- `movlw k` - load literal value
- `movwf MYVAR` - move value but do not affect status bits
- `movff V1,V2` - move data at source address to data at destination address
- `movf f,F/W` - move value and affect status bits
 - What does `movf COUNT,F` do?

Other move/load instructions

- `lrsr 0,num1` - load FSR register
 - 1st argument is FSR register number
 - 2nd argument is value to load
- Saving a FSR register
 - `movff FSR0L,FSR0L_TEMP`
 - `movff FSR0H,FRS0H_TEMP`
- Loading a FSR register
 - `movff FSR0L_TEMP,FSR0L`
 - `movff FRS0H_TEMP,FSR0H`

Load BSR register & other

- `movlb 2`
 - Load the value of 2 into the BSR register
- `clrf TEMP` – Load 0x00 into variable TEMP
- `setf TEMP` – Load 0xff into variable TEMP
- `swapf f` - swap nibbles of f

Single operand instructions

- Single bit instructions
 - `bsf PORTB,0` - Set the lsb of PORTB
 - `bcl PORTB,1` - clear bit 1 of PORTB
 - `btg PORTB,2` - toggle bit 2 of PORTB
- Rotate instructions illustrated on next slide
 - `rlcf rlncl rrcf rrcncf`
 - `cf` rotate including carry bit
 - `ncf` rotate without carry bit

Logical instructions

- `andlw B'00001111'` And WREG with value
- `andwf f,F/W` - AND WREG with f putting result back in F or WREG
- `iorlw k` -Inclusive OR literal value and WREG
- `iorwf f,F/W` – inclusive OR WREG with f and put result into f or WREG
- `xorlw k, xorwf f,F/W` - Exclusive OR

Arithmetic

- `addlw k, addwf f,F/W`
- `addwfc f,F/W` - Add WREG, F and carry bit
- `daw` – Decimal adjust sum of pack BCD
- `subwf, sublw`
- `subwfb f,F/W` - subtract with borrow

Multiplicaiton

- `mullw k` - multiply WREG with literal value k putting result in PRODH:PRODL reg - WREG unaffected
- `mullwf f(,Banked)` - Multiply WREG with f putting results in PRODH:PRODL - both WREG and f remain unchanged

Branches

- Needed for program flow control
- Tests on status register
 - bc, bnc, bz, bnz, bn, bnn, bov, bnov
 - Use the c, x, n, and ov bits of the status register
- bra – branch always

Conditional Skip instructions

- Ten further instructions that test for a condition and skip over the next instruction if the condition is met.
 - Next instruction is typically a branch or rcall
 - Very useful at end of a loop
- Loop code
- `decfsz count,F ;Decrement and skip if zero`
- `bra top_of_loop`

Skip instructions

- `cpfseq f` - skip if $f = \text{WREG}$
- `cpfsgt f` - skip if $f > \text{WREG}$
- `cpfslt f` - skip if $f < \text{WREG}$
- `tstfsz t` - Test f , skip if zero
- `decfsz f,F/W` - Decr f , res->WREG, skip if 0
- `dcfsnz f,F/W` - Decr f , res->WREG, skip not 0
- `incfsz f,F/W` - Incr f , res->WREG, skip if 0
- `infsnz f,F/W` - Incr f , res->WREG, skip not 0

Other – Subroutine, interrupt

- rcall label - call subroutine (within 512 instr)
- call label – call subroutine (anywhere)
- call label, FAST - call subroutine, copy state to shadow registers
- return – return from subroutine
- return FAST - return and restore from shadow registers
- return k - return and put value k in WREG

cont

- retfie - return from interrupt and re-enable
- retfie FAST – return, re-enable- restore
- push - Push addr of next instruction onto stack
- pop - discard address on top of stack
- clrwdt - clear watchdog timer
- sleep - Go into standby mode
- reset - software controlled reset
- nop

Review: PIC instructions

- Logical operations
 - andlw/andwf
 - iorlw/iorwf
 - xorlw/xorwf
- Rotates
 - rrf
 - rlf
- Jumps/calls/return
 - goto
 - call
 - return/retlw/retfie
- Miscellaneous
 - nop
 - sleep/clrwdt

Conditional Execution

STATUS bits:

none

- Conditional execution in PIC: skip next instruction if condition true
- Two general forms
 - Test bit and skip if bit clear/set
 - Increment/decrement register and skip if result is 0

btfsc	f, b	;Test bit b of register f, where b=0 to 7, skip if clear
btfss	f, b	;Test bit b of register f, where b=0 to 7, skip if set
decfsz	f, F(W)	;decrement f, putting result in F or W, skip if zero
incfsz	f, F(W)	;increment f, putting result in F or W, skip if zero

Examples:

- **btfsc** TEMP1, 0 ; Skip the next instruction if bit 0 of TEMP1 equals 0
- **btfss** STATUS, C ; Skip the next instruction if C==1
- **decfsz** TEMP1, F ; Decrement TEMP1, skip if TEMP1==0
- **incfsz** TEMP1, W ; W <- TEMP1+1 , skip if W==0 (TEMP1==0xFF)
; Leave TEMP1 unchanged

Example

- Show the values of all changed registers after each of the following sequences
 - What high-level operation does each perform?

(a)

```
movf    a, W
sublw   0xA
btfsc   STATUS, Z
goto    L1
incf    b, W
goto    L2

L1
decf    b, W

L2
movwf   a
```

(b)

```
movf    NUM2, W
subwf   NUM1, W
btfss   STATUS, C
goto    BL
movf    NUM1, W
goto    Done

BL
movf    NUM2, W

Done
movwf   MAX
```

Example solution (part a)

movf	a, W	→ $W = a$
sublw	0xA	→ $W = 10 - a$
btsc	STATUS, Z	→ Skip goto if result is non-zero
goto	L1	→ Goto L1 if result == 0 → Reach this point if result non-zero
incf	b, W	→ $W = b + 1$
goto	L2	
decf	b, W	→ $W = b - 1$
movwf	a	→ $a = W$ → value depends on what's executed before this

High-level operation:
if $((10 - a) == 0)$
 $a = b - 1$
else
 $a = b + 1$

L1

L2

Example solution (part b)

movf	NUM2, W	→ $W = \text{NUM2}$
subwf	NUM1, W	→ $W = \text{NUM1} - W$ = $\text{NUM1} - \text{NUM2}$
btfss	STATUS, C	→ Carry indicates “above” → if set, $\text{NUM1} > \text{NUM2}$
goto	BL	
movf	NUM1, W	→ if ($\text{NUM1} \geq \text{NUM2}$) $W = \text{NUM1}$
goto	Done	→ Skip “below” section
BL		
movf	NUM2, W	→ if ($\text{NUM1} < \text{NUM2}$) $W = \text{NUM2}$
Done		
movwf	MAX	

High-level operation:
if ($\text{NUM1} < \text{NUM2}$)
$\text{MAX} = \text{NUM2}$
else
$\text{MAX} = \text{NUM1}$

Working with 16-bit data

Assume a 16-bit counter, the upper byte of the counter is called COUNTH and the lower byte is called COUNTL.

Decrement a 16-bit counter

```
movf    COUNTL, F           ; Set Z if lower byte == 0
btfsc   STATUS, Z           ; if not, then done testing
decf    COUNTH, F          ; if so, decrement COUNTH
decf    COUNTL, F          ; in either case decrement COUNTL
```

Test a 16-bit variable for zero

```
movf    COUNTL, F           ; Set Z if lower byte == 0
btfsc   STATUS, Z           ; if not, then done testing
movf    COUNTH, F          ; Set Z if upper byte == 0
btfsc   STATUS, Z           ; if not, then done
goto    BothZero           ; branch if 16-bit variable == 0
```

CarryOn

A Delay Subroutine

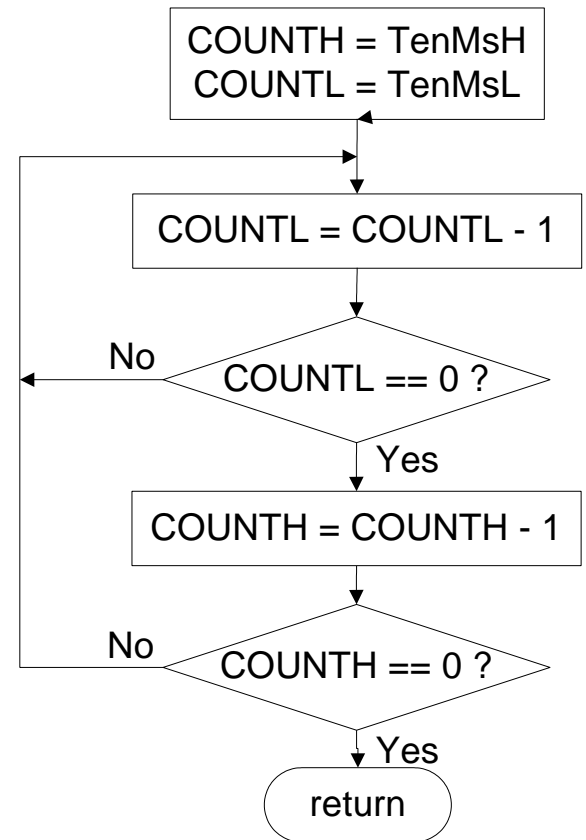
```
. *****  
;  
; TenMs subroutine and its call inserts a delay of exactly ten milliseconds  
; into the execution of code.  
; It assumes a 4 MHz crystal clock. One instruction cycle = 4 * Tosc.  
; TenMsH equ 13 ; Initial value of TenMs Subroutine's counter  
; TenMsL equ 250  
; COUNTH and COUNTL are two variables
```

TenMs

```
  nop ; one cycle  
  movlw TenMsH ; Initialize COUNT  
  movwf COUNTH  
  movlw TenMsL  
  movwf COUNTL
```

Ten_1

```
  decfsz COUNTL,F ; Inner loop  
  goto Ten_1  
  decfsz COUNTH,F ; Outer loop  
  goto Ten_1  
  return
```



Applications

- Personal information products: Cell phone, pager, watch, pocket recorder, calculator
- Laptop components: mouse, keyboard, modem, fax card, sound card, battery charger
- Home appliances: door lock, alarm clock, thermostat, air conditioner, tv remote, hair dryer, VCR, small refrigerator, exercise equipment, washer/dryer, microwave oven
- Toys; video games, cars, dolls, etc.

Summary

- Microprocessors and embedded controllers are a ubiquitous part of life today
- These devices come in a wide variety of configurations and designs
- Headhunters report that EEs familiar with μC , μP design are in the highest possible demand
- Feedback

Assignment

- Illustrate PIC instruction set with example.